# THE SELECTION OF PROCESS SEQUENCE FOR MULTI MACHINE ARRANGED IN SERIES IN THE JOB SHOP INDUSTRY 

Hery Irwan ${ }^{1,2^{*}}$, Md. Nizam Abd Rahman ${ }^{1}$, Zuhriah binti Ebrahim ${ }^{1}$, Tamara Handini ${ }^{2}$<br>${ }^{1}$ Faculty of Manufacturing Engineering, Universiti Teknikal Melaka, Malacca, Malaysia<br>${ }^{2}$ Faculty of Engineering, Universitas Riau Kepulauan, Batam, Indonesia<br>*Corresponding author: hery04@ gmail.com


#### Abstract

PT. KOP operates within the oil and gas drilling sector, utilizing job shop production scheduling. The company produces 4 primary product types, employing 2 identical machines operated in parallel, as is common in the job shop industry. However, production planning faces challenges, notably delivery delays stemming from inefficient scheduling, indicated by high work-in-process inventory and machine tooling availability issues. This research proposes a series-based machine scheduling method, evaluating dispatching methods such as SPT (Shortest Processing Time), LPT (Longest Processing Time), and the Johnson algorithm. Results demonstrate that the Johnson method outperforms SPT and LPT, minimizing delay times to 1 hour in the first week, 16 hours in the second, 48 hours in the third, and 86 hours in the fourth. By arranging machines in series, tooling requirements are reduced by $50 \%$. Specifically, series machines only require 4 tooling per product, while parallel machines require 8.


Keywords: Dispatching; LPT; Johnson Algorithm; Series; SPT

## INTRODUCTION

In today's competitive landscape, many industrial enterprises strive for exceptional outcomes to ensure customer satisfaction and minimize disappointment. Maintaining product quality is paramount for business success as it directly impacts customer happiness. Timely delivery also plays a crucial role in customer satisfaction. One of the significant challenges faced by organizations is imprecise production scheduling, resulting in frequent product delivery delays. Implementing an effective production scheduling procedure is essential to ensure timely order fulfillment and minimize these delays (Mansouri, Golmohammadi and Miller, 2019).

The "Scheduling" entails arranging and organizing the sequence of tasks or operations within a production line to be performed across multiple machines. Meanwhile, "processing time" refers to the duration required to execute an activity or operation, which involves allocating specific resources within a predefined timeframe (Golmohammadi and Mansouri, 2015). The challenge of scheduling in dynamic job shops revolves around creating schedules across multiple machines using dispatching rules. The choice of dispatching rule depends on the scheduling objective. Various time-based objectives are linked with scheduling (Holthaus and Rajendran, 2000; Hildebrandt, Heger and Scholz-Reiter, 2010; Branke, Hildebrandt and Scholz-Reiter, 2015).

PT. KOP operates in the oil and gas drilling sector and offers 4 distinct products. The company utilizes 2 machines arranged in parallel for production, with each product completing its manufacturing process across both machines with two setups. Unfortunately, PT. KOP frequently encounters delivery delays stemming from inadequate production scheduling. The company's production process involves a set of identical machines, each assigned to manufacture a specific product. Machine one is dedicated to completing one product, while machine two is responsible for producing another product. These machines operate independently and are not synchronized during the production process (Lin and Huang, 2021). Despite using two identical machines in the production process, the company often faces delays of one to two weeks in product delivery. This delay arises from the necessity to dispatch items in pairs, which is not fulfilled due to incomplete production. This inadequacy stems from improper implementation of the production scheduling technique. To rectify this issue, organizations must establish an appropriate dispatching rule to effectively manage scheduling (Mönch et al., 2005). In production scheduling, organizations often resort to random techniques, whereas in product delivery, they tend to operate in pairs. For example, product delivery companies may handle one to two products or three to four products simultaneously (Davari et al., 2020).

The aim of production scheduling is to reduce delay time from specified deadlines to meet customer demands. It also seeks to minimize the total processing time of all jobs and improve machine efficiency while reducing idle time (Grundstein, Freitag and Scholz-Reiter, 2017; Irwan, 2020). Dispatching rules comprise a set of guidelines used in scheduling operations to establish the sequence in which jobs are processed on individual machines. Priorities are assigned to each job in the queue to determine this order (Sharda, 2013).

The FCFS (First-Come, First-Served) dispatching rule is widely used in production scheduling and order management. It prioritizes orders or jobs based on their arrival time, ensuring that those received first are processed first. Conversely, the SPT (Shortest Processing Time) rule is another method employed in production scheduling. It prioritizes orders or jobs based on their processing times, giving precedence to those with the shortest duration. (Jayamohan and Rajendran, 2000; Shady et al., 2021).

The Johnson method is a technique in production and operations management employed to identify the most efficient sequence of tasks for executing the production process (Pinedo, 2008; Irwan, 2022). A Gantt chart is a type of bar chart that visually represents tasks
and their corresponding timeframes. It is commonly used for project planning and coordinating various scheduled activities (Baker and Scudder, 1990; Supithak, Liman and Montes, 2010).

Research conducted by Waschneck et al. (2017) and Pfund, Mason and Fowler (2006), in the semiconductor industry highlights the continued use of dispatching rules for shop floor control in wafer manufacture. This system offers real-time capability and ease of comprehension for operators. Among the dispatching techniques, SPT stands out as particularly effective in reducing work-in-progress (WIP) on the machine. The research was conducted using the dispatching rule approach, specifically employing the SPT, LPT and the Johnson approach.

## METHODOLOGY

The study took place at PT. KOP, a company operating in the oil and gas drilling industry. It focused on the machining operations, where the production of a single product involves two setup procedures. The data collected from the research include: (1) Production process time; (2) Setup; and (3) Demand. Then following the data analyze process based on the results of previous data processing attempts. The data processing and analysis involve several steps as describes in Table 1.

Table 1. Data Processing and Analysis Step

| Steps |
| :---: |
| Data processing |
| Gantt Chart Generating |
| Processing the data to establish a production <br> scheduling sequence using the Shortest <br> Processing Time (SPT), Longest Processing Time <br> (LPT), and Johnson methods to determine <br> delivery times |
| Select suitable method |
| Generating a Gantt chart from the data obtained <br> using the three methods. This chart serves as a <br> planning tool for resource scheduling and time <br> allocation |
| Identifying the most suitable dispatching rules <br> among the three methods (SPT, LPT, and <br> Johnson) for scheduling at PT. KOP and <br> understanding tooling utilization |

$1^{\text {st }}$ International Conference on Multidisciplinary Studies
Universitas Riau Kepulauan, Batam, December 19, 2023
e-ISSN: 3047-6399
Table 2 displays the order data for the month along with the corresponding due date. Furthermore, the manufacturing time and setup time are display in Table 3.

Table 2. Demand and Due Data (Hours)

| Data |  | Week 1 |  | Week 2 |  | Week 3 |  | Week 4 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Demand | Due | Demand | Due | Demand | Due | Demand | Due |
| Product 1 | Cust\# 1 |  | 144 | 2 | 288 |  | 432 | 2 | 576 |
|  | Cust\# 2 | 2 |  |  |  |  |  |  |  |
|  | Cust\# 3 |  |  | 4 |  | 4 |  | 4 |  |
|  | Cust\# 4 | 2 |  |  |  | 2 |  |  |  |
|  | Total | 4 | 144 | 6 | 288 | 6 | 432 | 6 | 576 |
| Product 2 | Cust\# 1 |  | 144 | 2 | 288 |  | 432 | 3 | 576 |
|  | Cust\# 2 | 4 |  |  |  |  |  |  |  |
|  | Cust\# 3 |  |  | 4 |  | 5 |  | 4 |  |
|  | Cust\# 4 | 4 |  |  |  | 2 |  |  |  |
|  | Total | 8 | 144 | 6 | 288 | 7 | 432 | 7 | 576 |
| Product 3 | Cust\# 1 | 1 | 144 |  | 288 | 2 | 432 | 1 | 576 |
|  | Cust\# 2 |  |  | 2 |  | 1 |  | 1 |  |
|  | Cust\# 3 | 1 |  | 2 |  | 1 |  |  |  |
|  | Cust\# 4 | 1 |  |  |  |  |  | 2 |  |
|  | Total | 3 | 144 | 5 | 288 | 4 | 432 | 4 | 576 |
| Product 4 | Cust\# 1 | , | 144 |  | 288 | 2 | 432 | 1 | 576 |
|  | Cust\# 2 |  |  | 1 |  | 1 |  | 1 |  |
|  | Cust\# 3 | 1 |  | 1 |  | 1 |  |  |  |
|  | Cust\# 4 | 1 |  | 1 |  |  |  |  |  |
|  | Total | 3 | 144 | 3 | 288 | 4 | 432 | 2 | 576 |

Table 3. Data Processing Time and Setup

| No | Item | 1st Process (Hours) |  | 2nd Process (Hours) |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Setup | Process | Setup | Process |
| 1 | Product 1 | 2 | 7 | 2 | 10 |
| 2 | Product 2 | 1 | 3 | 1 | 6 |
| 3 | Product 3 | 1 | 3 | 1 | 6 |
| 4 | Product 4 | 2 | 8 | 2 | 11 |

## Current Schedule

The current schedule algorithm prioritizes the execution of tasks based on their arrival order, ensuring that the earliest arriving task is processed first. The production process adheres to the FCFS (First Come, First Served) rule, wherein the production sequence extends from the first week to the fourth week, aligning with the product sequence 1-2-3-4. The computation of flow time is illustrated in Table 4.

## Proceeding

$1^{\text {st }}$ International Conference on Multidisciplinary Studies
Universitas Riau Kepulauan, Batam, December 19, 2023
e-ISSN: 3047-6399
Table 4. Data Current Schedule - Job Shop

| Machine | Week 1 |  | Week 2 |  | Week 3 |  | Week 4 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Seq\# | Plan | Seq\# | Plan | Seq\# | Plan | Seq\# | Plan |
| $\#$ | 1 | 4 | 1 | 6 | 1 | 6 | 1 | 10 |
|  | 2 | 8 | 2 | 6 | 2 | 7 | 2 | 6 |
| $\# 2$ | 3 | 3 | 3 | 5 | 3 | 4 | 3 | 6 |
|  | 4 | 3 | 4 | 3 | 4 | 4 | 4 | 11 |

The sequential arrangement of the manufacturing process for obtaining flow time data using FCFS scheduling policy is presented in Table 5.

Table 5. Flow Time Machine Schedule using FCFS

| M | Week 1 (144) |  |  |  | Week 2 (288) |  |  |  | Week 3 (342) |  |  |  | Week 4 (576) |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Seq\# | Req | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req | $\mathrm{F}_{\mathrm{t}}$ | Late |
| \#1 | $1_{1}$ | 30 | 30 | 0 | $1_{1}$ | 44 | 190 | 0 | $1_{1}$ | 44 | 352 | 0 | $1_{1}$ | 44 | 523 | 0 |
|  | $1_{2}$ | 42 | 72 | 0 | $1_{2}$ | 62 | 252 | 0 | $1_{2}$ | 62 | 414 | 0 | $1_{2}$ | 62 | 585 | 9 |
|  | 21 | 25 | 97 | 0 | 21 | 19 | 271 | 0 | 21 | 22 | 436 | 4 | 21 | 22 | 607 | 31 |
|  | $2{ }_{2}$ | 69 | 146 | 2 | $2{ }_{2}$ | 37 | 308 | 20 | $2{ }_{2}$ | 43 | 479 | 47 | $2{ }_{2}$ | 43 | 650 | 74 |
| \#2 | $3_{1}$ | 10 | 10 | 0 | $3_{1}$ | 16 | 106 | 0 | $3_{1}$ | 13 | 211 | 0 | $3_{1}$ | 13 | 329 | 0 |
|  | $3_{2}$ | 19 | 29 | 0 | 32 | 31 | 137 | 0 | $3_{2}$ | 25 | 236 | 0 | 32 | 25 | 354 | 0 |
|  | 41 | 26 | 55 | 0 | 41 | 26 | 163 | 0 | $4_{1}$ | 34 | 270 | 0 | 41 | 34 | 388 | 0 |
|  | 42 | 35 | 90 | 0 | 42 | 35 | 198 | 0 | 42 | 46 | 316 | 0 | 42 | 46 | 434 | 0 |

From the data provided in the table, it's clear that implementing the FCFS strategy in production planning at the company results in delays. In the first week, there's a 2-hour delay observed for product \#2 on machine \#1. In the second week, there's a delay of 20 hours for producing product \#2. Product \#2 encounters a delay of 51 hours in the third week. In the fourth week, both product \#1 and product \#2 experience delays, amounting to 9 hours and 105 hours respectively, compared to the scheduled timeframe.

## Shortest Processing Time (SPT)

Afterwards, machine scheduling will be conducted sequentially following a priority order determined by the SPT (Shortest Processing Time) method. Scheduling jobs based on SPT calculations prioritizes those with the shortest production processing time for completion. The following table illustrates the computation of the scheduling order using the SPT rule. Table 6 demonstrates the sequence of processes for the demand over a four-week period, employing the SPT approach.

Table 6. Prioritize Planning using SPT

| Item | Sub Item | Week 1 |  | Week 2 |  | Week 3 |  | Week 4 |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\mathrm{P}_{\mathrm{t}}$ | Total | $\mathrm{P}_{\mathrm{t}}$ | Total | $\mathrm{P}_{\mathrm{t}}$ | Total | $\mathrm{P}_{\mathrm{t}}$ | Total |
| Prod \#1 | $1_{1}$ | 30 | $72^{(\mathbf{3})}$ | 44 | $106^{(\mathbf{4})}$ | 44 | $106^{(\mathbf{4})}$ | 44 | $106^{(\mathbf{4})}$ |
|  | $1_{2}$ | 42 |  | 62 |  | 62 |  | 62 |  |
| Prod \#2 | $2_{1}$ | 25 | $74^{\mathbf{( 4 )}}$ | 19 | $56^{(\mathbf{2})}$ | 22 | $65^{(\mathbf{2})}$ | 22 | $65^{(\mathbf{2})}$ |
|  | $2_{2}$ | 49 |  | 37 |  | 43 |  | 43 |  |
| Prod \#3 | $3_{1}$ | 10 | $29^{(\mathbf{1})}$ | 16 | $47^{(\mathbf{1})}$ | 13 | $38^{(\mathbf{1})}$ | 13 | $38^{(\mathbf{1})}$ |
|  | $3_{2}$ | 19 |  | 31 |  | 25 |  | 25 |  |
| Prod \#4 | $4_{1}$ | 26 | $61^{\mathbf{( 2 )}}$ | 26 | $61^{\mathbf{( 3 )}}$ | 34 | $80^{(\mathbf{3})}$ | 34 | $80^{(\mathbf{3})}$ |
|  | $4_{2}$ | 35 |  | 35 |  | 46 |  | 46 |  |

The results of the preceding calculations can be arranged in accordance with the SPT regulations, the order calculation is performed each week as follows:
a) First-week product scheduling sequence 3-2-1-4
b) Second-week product scheduling sequence 3-2-4-1
c) Third week product scheduling sequence 3-2-4-1
d) Product 3-2-4-1 scheduling sequence for the fourth week.

Data flow time based on SPT principles, as shown in Table 7.
Table 7. Flow Time Machine Schedule using SPT

| M | Week 1 (144) |  |  |  | Week 2 (288) |  |  |  | Week 3 (432) |  |  |  | Week 4 (576) |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Seq | Req | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Lat | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Late |
| \#1 | 31 | 10 | 10 | 0 | 31 | 16 | 107 | 0 | 31 | 13 | 209 | 0 | 31 | 13 | 322 | 0 |
|  | 21 | 25 | 35 | 0 | 21 | 19 | 126 | 0 | 21 | 22 | 231 | 0 | 21 | 22 | 344 | 0 |
|  | $1_{1}$ | 30 | 65 | 0 | 41 | 26 | 152 | 0 | 41 | 34 | 265 | 0 | 41 | 34 | 378 | 0 |
|  | 41 | 26 | 91 | 0 | $1_{1}$ | 44 | 196 | 0 | $1_{1}$ | 44 | 309 | 0 | $1_{1}$ | 44 | 422 | 0 |
| \#2 | 32 | 19 | 19 | 0 | 31 | 16 | 161 | 0 | 32 | 25 | 320 | 0 | 32 | 25 | 496 | 0 |
|  | 22 | 49 | 68 | 0 | $2{ }_{2}$ | 37 | 198 | 0 | $2{ }_{2}$ | 43 | 363 | 0 | $2{ }_{2}$ | 43 | 539 | 0 |
|  | 12 | 42 | 110 | 0 | 42 | 35 | 233 | 0 | 42 | 46 | 409 | 0 | 42 | 46 | 585 | 9 |
|  | 42 | 35 | 145 | 1 | $1_{2}$ | 62 | 295 | 7 | $1_{2}$ | 62 | 471 | 39 | $1_{2}$ | 62 | 647 | 71 |

Based on the data presented in the table, it is evident that the implementation of the SPT approach for production planning throughout the organization leads to delays specifically on machine \#2. In the first week delay for 1 hour on last product, the second week there was a delay of 7 hours for product \#4. In the third week, a delay of 39 hours has been observed to produce product \#4. In the fourth week, there is a delay of 9 hours for product \#2 and 71 hours for product \#4 from capacity on the week limits.

## Proceeding

$I^{\text {st }}$ International Conference on Multidisciplinary Studies
Universitas Riau Kepulauan, Batam, December 19, 2023
e-ISSN: 3047-6399

## Longest Processing Time (LPT)

Afterwards, machine scheduling will be conducted sequentially following a priority order determined by the LPT (Longest Processing Time) method. Scheduling jobs based on LPT calculations prioritizes those with the longest production processing time for completion. The following table illustrates the computation of the scheduling order using the LPT rule. Table 8 demonstrates the sequence of processes for the demand over a four-week period, employing the LPT approach.

Table 8. Prioritize Planning using LPT

| Item | Sub Item | Week 1 |  | Week 2 |  | Week 3 |  | Week 4 |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\mathrm{P}_{\mathrm{t}}$ | Total | $\mathrm{P}_{\mathrm{t}}$ | Total | $\mathrm{P}_{\mathrm{t}}$ | Total | $\mathrm{P}_{\mathrm{t}}$ | Total |
| Prod \#1 | $1_{1}$ | 30 | $72^{(\mathbf{1 )}}$ | 44 | $106^{(\mathbf{1 )}}$ | 44 | $106^{(1)}$ | 44 | $106^{(\mathbf{1 )}}$ |
|  | $1_{2}$ | 42 |  | 62 |  | 62 |  | 62 |  |
| Prod \#2 | $2_{1}$ | 25 | $74^{(3)}$ | 19 | $56^{(3)}$ | 22 | $65^{(3)}$ | 22 | $65^{(3)}$ |
|  | $2_{2}$ | 49 |  | 37 |  | 43 |  | 43 |  |
| Prod \#3 | $3_{1}$ | 10 | $29^{(4)}$ | 16 | $47^{(4)}$ | 13 | $38^{(4)}$ | 13 | $38^{(4)}$ |
|  | $3_{2}$ | 19 |  | 31 |  | 25 |  | 25 |  |
| Prod \#4 | $4_{1}$ | 26 | $61^{(2)}$ | 26 | $61^{(2)}$ | 34 | $80^{(2)}$ | 34 | $80^{(2)}$ |
|  | $4_{2}$ | 35 |  | 35 |  | 46 |  | 46 |  |

The results of the preceding calculations can be arranged in accordance with the SPT regulations, the order calculation is performed each week as follows.
e) First-week product scheduling sequence 1-4-2-3
f) Second-week product scheduling sequence 1-4-2-3
g) Third week product scheduling sequence 1-4-2-3
h) Product 1-4-2-3scheduling sequence for the fourth week.

Data flow time based on LPT principles, as shown in Table 9.
Table 9. Flow Time Machine Schedule using LPT

| M | Week 1 (144) |  |  |  | Week 2 (288) |  |  |  | Week 3 (432) |  |  |  | Week 4 (576) |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Seq \# | Req | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | $\begin{gathered} \hline \text { Lat } \\ \mathrm{e} \end{gathered}$ | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Late |
| \#1 | $1{ }_{1}$ | 30 | 30 | 0 | $1_{1}$ | 44 | 74 | 0 | $1_{1}$ | 44 | 179 | 0 | $1_{1}$ | 44 | 292 | 0 |
|  | 4 | 26 | 56 | 0 | $4_{1}$ | 26 | 100 | 0 | $4_{1}$ | 34 | 213 | 0 | $4_{1}$ | 34 | 326 | 0 |
|  | 21 | 25 | 81 | 0 | 21 | 19 | 119 | 0 | 21 | 22 | 235 | 0 | 21 | 22 | 348 | 0 |
|  | $3_{1}$ | 10 | 91 | 0 | $3_{1}$ | 16 | 135 | 0 | $3_{1}$ | 13 | 248 | 0 | $3_{1}$ | 13 | 361 | 0 |
| \#2 | $1_{2}$ | 42 | 42 | 0 | 12 | 62 | 207 | 0 | $1_{2}$ | 62 | 357 | 0 | 12 | 62 | 533 | 0 |
|  | 42 | 35 | 77 | 0 | 42 | 35 | 242 | 0 | 42 | 46 | 403 | 0 | 42 | 46 | 579 | 3 |
|  | 22 | 49 | 126 | 0 | $2{ }_{2}$ | 37 | 279 | 0 | $2{ }_{2}$ | 43 | 446 | 14 | $2{ }_{2}$ | 43 | 622 | 46 |
|  | 32 | 19 | 145 | 1 | 31 | 16 | 295 | 7 | 32 | 25 | 471 | 39 | 32 | 25 | 647 | 71 |

Based on the data presented in the table, it is evident that the implementation of the
LPT approach for production planning throughout the organization leads to delays
specifically on machine \#2. In the first week delay for 1 hour on last product, the second week there was a delay of 7 hours for product \#3. In the third week, a delay of 39 hours has been observed to produce product \#3 and 14 hours for product \#2. In the fourth week, there is a delay happen on 3 product there is 3 hours for product \#4, 46 hours for product \#2 and 71 hours for product \#3 from capacity on the week limits.

## Johson Algorithm

The Johnson method entails computing the processing time required for each operation across all jobs. The aim is to determine the operation with the shortest processing duration. If the shortest operation occurs at the beginning of a job, it should be performed on the last machine in the production sequence. Conversely, if the shortest operation occurs at the end of a job, it should be executed on the first machine in the production sequence. Initially, the shortest operation is removed from the first machine in the production sequence. Then, the operation with the shortest processing time at that moment is identified. This process iterates until all operations have been completed.

Johnson algorithm for two-stage production schedule as follow: Divide the job's set into two disjoint subsets, $\operatorname{Sub}_{1}$ and $\operatorname{Sub}_{2} ;$ where $\operatorname{Sub}_{1}=\left\{\mathrm{J}_{\mathrm{i}}, \mathrm{p}_{\mathrm{i} 1} \leq \mathrm{p}_{\mathrm{i} 2}\right\}$ and $\operatorname{Sub}_{2}=\left\{\mathrm{J}_{\mathrm{i}}, \mathrm{p}_{\mathrm{i} 1}>\mathrm{p}_{\mathrm{i} 2}\right\}$. Order the jobs within $\mathrm{Sub}_{1}$ in the non-decreasing order of $\mathrm{p}_{\mathrm{i} 1}$ and those in $\mathrm{Sub}_{2}$ in the nonincreasing order of $\mathrm{p}_{\mathrm{i} 2}$. Sequence jobs in $\mathrm{Sub}_{1}$ first, followed by $\mathrm{Sub}_{2}$ (Benttaleb, Hnaien and Yalaoui, 2018). Please refer to Table 10 for specifics on the procedure for determining the sequence using the Johnson algorithm.

Table 10. Prioritize Planning using Johnson Algorithm

| Item | Week 1 |  | Week 2 |  | Week 3 |  | Week 4 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Proc \#1 | Proc \#2 | Proc \#1 | Proc \#2 | Proc \#1 | Proc \#2 | Proc \#1 | Proc \#2 |
| Product 1 | $30^{(4)}$ | 42 | $54^{(3)}$ | 66 | $78{ }^{(3)}$ | 90 | $102{ }^{(3)}$ | 114 |
| Product 2 | $25^{(2)}$ | 49 | $73{ }^{(4)}$ | 97 | $121{ }^{(4)}$ | 145 | $169{ }^{(4)}$ | 193 |
| Product 3 | $10^{(1)}$ | 19 | $28^{(1)}$ | 37 | $46^{(1)}$ | 55 | $64{ }^{(1)}$ | 73 |
| Product 4 | $26^{(3)}$ | 35 | $44^{(2)}$ | 53 | $62^{(2)}$ | 71 | $80^{(2)}$ | 89 |
| Sequences | 3-2-4-1 |  | 3-4-1-2 |  | 3-4-1-2 |  | 3-4-1-2 |  |

The results of the preceding calculations can be arranged in accordance with the Johnson Algorithm, the order calculation is performed each week as follows.
a) First-week product scheduling sequence 3-2-4-1
b) Second-week product scheduling sequence 3-4-1-2
c) Third week product scheduling sequence 3-4-1-2

Ist $^{\text {st }}$ International Conference on Multidisciplinary Studies
Universitas Riau Kepulauan, Batam, December 19, 2023
e-ISSN: 3047-6399
d) Product 3-4-1-2 scheduling sequence for the fourth week.

Data flow time based on Johnson Algorithm, as shown in Table 11.
Table 11. Flow Time Machine Schedule using Johnson Algorithm

| M | Week 1 (144) |  |  |  | Week 2 (288) |  |  |  | Week 3 (432) |  |  |  | Week 4 (576) |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Seq | Req | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Lat | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\# | Req. | $\mathrm{F}_{\mathrm{t}}$ | Late |
| \#1 | 31 | 10 | 10 | 0 | 31 | 16 | 26 | 0 | 31 | 13 | 128 | 0 | 31 | 13 | 241 | 0 |
|  | 21 | 25 | 35 | 0 | 41 | 26 | 52 | 0 | $4_{1}$ | 34 | 162 | 0 | 41 | 34 | 275 | 0 |
|  | 4 | 26 | 61 | 0 | $1_{1}$ | 44 | 96 | 0 | $1_{1}$ | 44 | 206 | 0 | $1{ }_{1}$ | 44 | 319 | 0 |
|  | $1_{1}$ | 30 | 91 | 2 | 21 | 19 | 115 | 0 | 21 | 22 | 228 | 0 | 21 | 22 | 341 | 0 |
| \#2 | 32 | 19 | 19 | 0 | 32 | 31 | 176 | 0 | 32 | 25 | 335 | 0 | 32 | 25 | 511 | 0 |
|  | 22 | 49 | 68 | 0 | $4{ }_{2}$ | 35 | 211 | 0 | 42 | 46 | 381 | 0 | $4{ }_{2}$ | 46 | 557 | 9 |
|  | 42 | 35 | 103 | 0 | $1_{2}$ | 62 | 273 | 0 | 12 | 62 | 443 | 11 | 12 | 62 | 619 | 43 |
|  | $1_{2}$ | 42 | 145 | 1 | $2{ }_{2}$ | 37 | 310 | 22 | $2{ }_{2}$ | 43 | 486 | 54 | 2 | 43 | 662 | 86 |

Based on the data presented in the table, it is evident that the implementation of the Johnson algorithm approach for production planning throughout the organization leads to delays specifically on machine \#2. In the first week delay for 1 hour on last product, the second week there was a delay of 22 hours for product \#1. In the third week, a delay of 11 hours has been observed to produce product \#4 and 54 hours for product \#1. In the fourth week, there is a delay happen on product \#4 about 43 hours and 86 hours on product \#1.

## RESULTS AND DISCUSSION

Based on the research results from the trial with three approaches incorporating completion time, work tardiness, and mean flow time calculations, it is discovered that SPT makes a greater contribution. Table 12 presents comprehensive data regarding delays that arise from the three dispatching techniques. This data is obtained by comparing the flowtime of each product with its respective due date upon completion on each machine. Delays are observed when the duration of the flowtime beyond the pre-established deadline for each individual product.

In Table 12, it is evident that the SPT method for scheduling machine \#2 results in a 1-hour delay for product \#4 in the first week. This delay extends into the second week, reaching 7 hours and impact to Product \#1, and by the third week still happen on product \#1 about 39 hours behind schedule. Finally, at the conclusion of the fourth week, there is 2 product impact for late delivery, that is product \#4 about 9 hours behind schedule and product \#1 about 71 hours behind the schedule. The SPT scheduling still better than another 2 model and current practice schedules done by Planners.

Ist $^{\text {st }}$ International Conference on Multidisciplinary Studies
Universitas Riau Kepulauan，Batam，December 19， 2023
e－ISSN：3047－6399
Table 12．Comparison of Tardiness Data

|  | M | Week 1 （144） |  |  |  | Week 2 （288） |  |  |  | Week 3 （342） |  |  |  | Week 4 （576） |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Seq \# | Req | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\＃ | Req． | $\mathrm{F}_{\mathrm{t}}$ | $\begin{gathered} \text { Lat } \\ \mathrm{e} \end{gathered}$ | Seq\＃ | Req． | $\mathrm{F}_{\mathrm{t}}$ | Late | Seq\＃ | Req． | $\mathrm{F}_{\mathrm{t}}$ | Late |
| 恶 | \＃1 | $1_{1}$ | 30 | 30 | 0 | $1_{1}$ | 44 | 190 | 0 | $1_{1}$ | 44 | 352 | 0 | $1_{1}$ | 44 | 523 | 0 |
|  |  | $1_{2}$ | 42 | 72 | 0 | 12 | 62 | 252 | 0 | 12 | 62 | 414 | 0 | 12 | 62 | 585 | 9 |
|  |  | 21 | 25 | 97 | 0 | 21 | 19 | 271 | 0 | 21 | 22 | 436 | 4 | 21 | 22 | 607 | 31 |
|  |  | 2 | 69 | 146 | 2 | 2 | 37 | 308 | 20 | 2 | 43 | 479 | 47 | 2 | 43 | 650 | 74 |
|  | \＃2 | 31 | 10 | 10 | 0 | 31 | 16 | 106 | 0 | 31 | 13 | 211 | 0 | 31 | 13 | 329 | 0 |
|  |  | $3_{2}$ | 19 | 29 | 0 | 32 | 31 | 137 | 0 | 32 | 25 | 236 | 0 | $3{ }_{2}$ | 25 | 354 | 0 |
|  |  | 41 | 26 | 55 | 0 | 41 | 26 | 163 | 0 | 41 | 34 | 270 | 0 | 41 | 34 | 388 | 0 |
|  |  | $4_{2}$ | 35 | 90 | 0 | 42 | 35 | 198 | 0 | $4_{2}$ | 46 | 316 | 0 | 42 | 46 | 434 | 0 |
| 㐫 | \＃1 | 31 | 10 | 10 | 0 | 31 | 16 | 107 | 0 | 31 | 13 | 209 | 0 | 31 | 13 | 322 | 0 |
|  |  | 21 | 25 | 35 | 0 | 21 | 19 | 126 | 0 | 21 | 22 | 231 | 0 | 21 | 22 | 344 | 0 |
|  |  | $1_{1}$ | 30 | 65 | 0 | 41 | 26 | 152 | 0 | 41 | 34 | 265 | 0 | 41 | 34 | 378 | 0 |
|  |  | $4_{1}$ | 26 | 91 | 0 | $1_{1}$ | 44 | 196 | 0 | $1_{1}$ | 44 | 309 | 0 | $1_{1}$ | 44 | 422 | 0 |
|  | \＃2 | 32 | 19 | 19 | 0 | 31 | 16 | 161 | 0 | $3_{2}$ | 25 | 320 | 0 | 32 | 25 | 496 | 0 |
|  |  | $2{ }_{2}$ | 49 | 68 | 0 | 22 | 37 | 198 | 0 | $2{ }_{2}$ | 43 | 363 | 0 | $2{ }_{2}$ | 43 | 539 | 0 |
|  |  | $1_{2}$ | 42 | 110 | 0 | 42 | 35 | 233 | 0 | 42 | 46 | 409 | 0 | 42 | 46 | 585 | 9 |
|  |  | 42 | 35 | 145 | 1 | $1_{2}$ | 62 | 295 | 7 | $1_{2}$ | 62 | 471 | 39 | $1_{2}$ | 62 | 647 | 71 |
| 気 | \＃1 | 1 | 30 | 30 | 0 | $1{ }_{1}$ | 44 | 74 | 0 | $1_{1}$ | 44 | 179 | 0 | $1{ }_{1}$ | 44 | 292 | 0 |
|  |  | 41 | 26 | 56 | 0 | 41 | 26 | 100 | 0 | 41 | 34 | 213 | 0 | 41 | 34 | 326 | 0 |
|  |  | 21 | 25 | 81 | 0 | 21 | 19 | 119 | 0 | 21 | 22 | 235 | 0 | 21 | 22 | 348 | 0 |
|  |  | 31 | 10 | 91 | 0 | 31 | 16 | 135 | 0 | 31 | 13 | 248 | 0 | 31 | 13 | 361 | 0 |
|  | \＃2 | $1_{2}$ | 42 | 42 | 0 | $1_{2}$ | 62 | 207 | 0 | $1_{2}$ | 62 | 357 | 0 | $1_{2}$ | 62 | 533 | 0 |
|  |  | $4_{2}$ | 35 | 77 | 0 | 42 | 35 | 242 | 0 | 42 | 46 | 403 | 0 | 42 | 46 | 579 | 3 |
|  |  | $2_{2}$ | 49 | 126 | 0 | $2{ }_{2}$ | 37 | 279 | 0 | $2{ }_{2}$ | 43 | 446 | 14 | $2{ }_{2}$ | 43 | 622 | 46 |
|  |  | $3_{2}$ | 19 | 145 | 1 | $3_{1}$ | 16 | 295 | 7 | $3_{2}$ | 25 | 471 | 39 | $3_{2}$ | 25 | 647 | 71 |
| E0弟－ | \＃1 | 31 | 10 | 10 | 0 | 31 | 16 | 26 | 0 | $3_{1}$ | 13 | 128 | 0 | 31 | 13 | 241 | 0 |
|  |  | 21 | 25 | 35 | 0 | 41 | 26 | 52 | 0 | 41 | 34 | 162 | 0 | 41 | 34 | 275 | 0 |
|  |  | 41 | 26 | 61 | 0 | $1_{1}$ | 44 | 96 | 0 | $1_{1}$ | 44 | 206 | 0 | $1_{1}$ | 44 | 319 | 0 |
|  |  | $1_{1}$ | 30 | 91 | 2 | 21 | 19 | 115 | 0 | 21 | 22 | 228 | 0 | 21 | 22 | 341 | 0 |
|  | \＃2 | 32 | 19 | 19 | 0 | $3{ }_{2}$ | 31 | 176 | 0 | 32 | 25 | 335 | 0 | 32 | 25 | 511 | 0 |
|  |  | 22 | 49 | 68 | 0 | 42 | 35 | 211 | 0 | $4{ }_{2}$ | 46 | 381 | 0 | 42 | 46 | 557 | 9 |
|  |  | 42 | 35 | 103 | 0 | $1_{2}$ | 62 | 273 | 0 | $1_{2}$ | 62 | 443 | 11 | $1_{2}$ | 62 | 619 | 43 |
|  |  | 12 | 42 | 145 | 1 | 22 | 37 | 310 | 22 | 22 | 43 | 486 | 54 | 22 | 43 | 662 | 86 |

## CONCLUSIONS

The findings from the research conducted on three dispatching strategies for scheduling a job shop when machine set as series has been give positive contribution on tardiness．Based on research for 2 machine identical set as series a method selection for priority sequence is SPT．That may still can be improve further when research more deep using optimization through simulation．

## REFERENCES

Baker，K．R．and Scudder，G．D．（1990）＇Sequencing with earliness and tardiness penalties．A review＇，Operations Research，38（1），pp．22－36．doi：10．1287／opre．38．1．22．

Benttaleb, M., Hnaien, F. and Yalaoui, F. (2018) 'Two-machine job shop problem under availability constraints on one machine: Makespan minimization', Computers and Industrial Engineering, 117, pp. 138-151. doi: 10.1016/j.cie.2018.01.028.

Branke, J., Hildebrandt, T. and Scholz-Reiter, B. (2015) 'Hyper-heuristic evolution of dispatching rules: A comparison of rule representations', Evolutionary Computation, 23(2), pp. 249-277. doi: 10.1162/EVCO_a_00131.

Davari, M. et al. (2020) 'Minimizing makespan on a single machine with release dates and inventory constraints', European Journal of Operational Research, 286(1), pp. 115128. doi: 10.1016/j.ejor.2020.03.029.

Golmohammadi, D. and Mansouri, S. A. (2015) 'Complexity and workload considerations in product mix decisions under the theory of constraints', Naval Research Logistics, 62(5), pp. 357-369. doi: 10.1002/nav. 21632.

Grundstein, S., Freitag, M. and Scholz-Reiter, B. (2017) 'A new method for autonomous control of complex job shops - Integrating order release, sequencing and capacity control to meet due dates', Journal of Manufacturing Systems, 42, pp. 11-28. doi: 10.1016/j.jmsy.2016.10.006.

Hildebrandt, T., Heger, J. and Scholz-Reiter, B. (2010) 'Towards improved dispatching rules for complex shop floor scenarios', Proceedings of the 12th annual conference on Genetic and evolutionary computation - GECCO '10, p. 257.

Holthaus, O. and Rajendran, C. (2000) 'Efficient jobshop dispatching rules: Further developments', Production Planning and Control, 11(2), pp. 171-178. doi: 10.1080/095372800232379.

Irwan, H. (2020) 'A Study Review of Completion Multi Job on Job Shop Scheduling Technique To Minimize Make Span', PROFISIENSI: Jurnal Program Studi Teknik Industri, 8(1), pp. 7-14. doi: 10.33373/profis.v8i1.2554.

Irwan, H. (2022) 'A review of integration model of lot-sizing-scheduling problem', Malaysian Construction Research Journal, 17(3), pp. 160-174. Available at: https://www.cream.my/usr/product.aspx?pgid=88\&id=58\&lang=en.

Jayamohan, M. S. and Rajendran, C. (2000) 'New dispatching rules for shop scheduling: A step forward', International Journal of Production Research, 38(3), pp. 563-586. doi: 10.1080/002075400189301.

Lin, D.-Y. and Huang, T.-Y. (2021) 'A Hybrid Metaheuristic for the Unrelated Parallel Machine Scheduling Problem', Mathematics, 9(7), p. 768. doi: 10.3390/math9070768.

Mansouri, S. A., Golmohammadi, D. and Miller, J. (2019) 'The moderating role of master production scheduling method on throughput in job shop systems', International Journal of Production Economics, 216(April), pp. 67-80.

Mönch, L. et al. (2005) 'Heuristic scheduling of jobs on parallel batch machines with incompatible job families and unequal ready times', Computers and Operations Research, 32(11), pp. 2731-2750. doi: 10.1016/j.cor.2004.04.001.

Pfund, M. E., Mason, S. J. and Fowler, J. W. (2006) 'Semiconductor Manufacturing Scheduling and Dispatching', Handbook of Production Scheduling, pp. 213-241. doi: 10.1007/0-387-33117-4_9.

Pinedo, M. L. (2008) Scheduling: Theory, algorithms, and systems, Scheduling: Theory, Algorithms, and Systems. Springer New York. doi: 10.1007/978-0-387-78935-4.

Shady, S. et al. (2021) 'Evolving Dispatching Rules Using Genetic Programming for Multiobjective Dynamic Job Shop Scheduling with Machine Breakdowns', in Procedia CIRP. Elsevier B.V., pp. 411-416. doi: 10.1016/j.procir.2021.11.069.

Sharda, R. (ed.) (2013) Production Planning and Control for Semiconductor Wafer Fabrication Facilities. Springer New York Heidelberg Dordrecht London. doi: 10.1007/978-1-4614-4472-5.

Supithak, W., Liman, S. D. and Montes, E. J. (2010) 'Lot-sizing and scheduling problem with earliness tardiness and setup penalties', Computers and Industrial Engineering, 58(3), pp. 363-372. doi: 10.1016/j.cie.2008.10.005.

Waschneck, B. et al. (2017) 'Production scheduling in complex job shops from an industrie 4.0 perspective: A review and challenges in the semiconductor industry', CEUR Workshop Proceedings, 1793.

